



Utility-Aware Social Event-Participant Planning

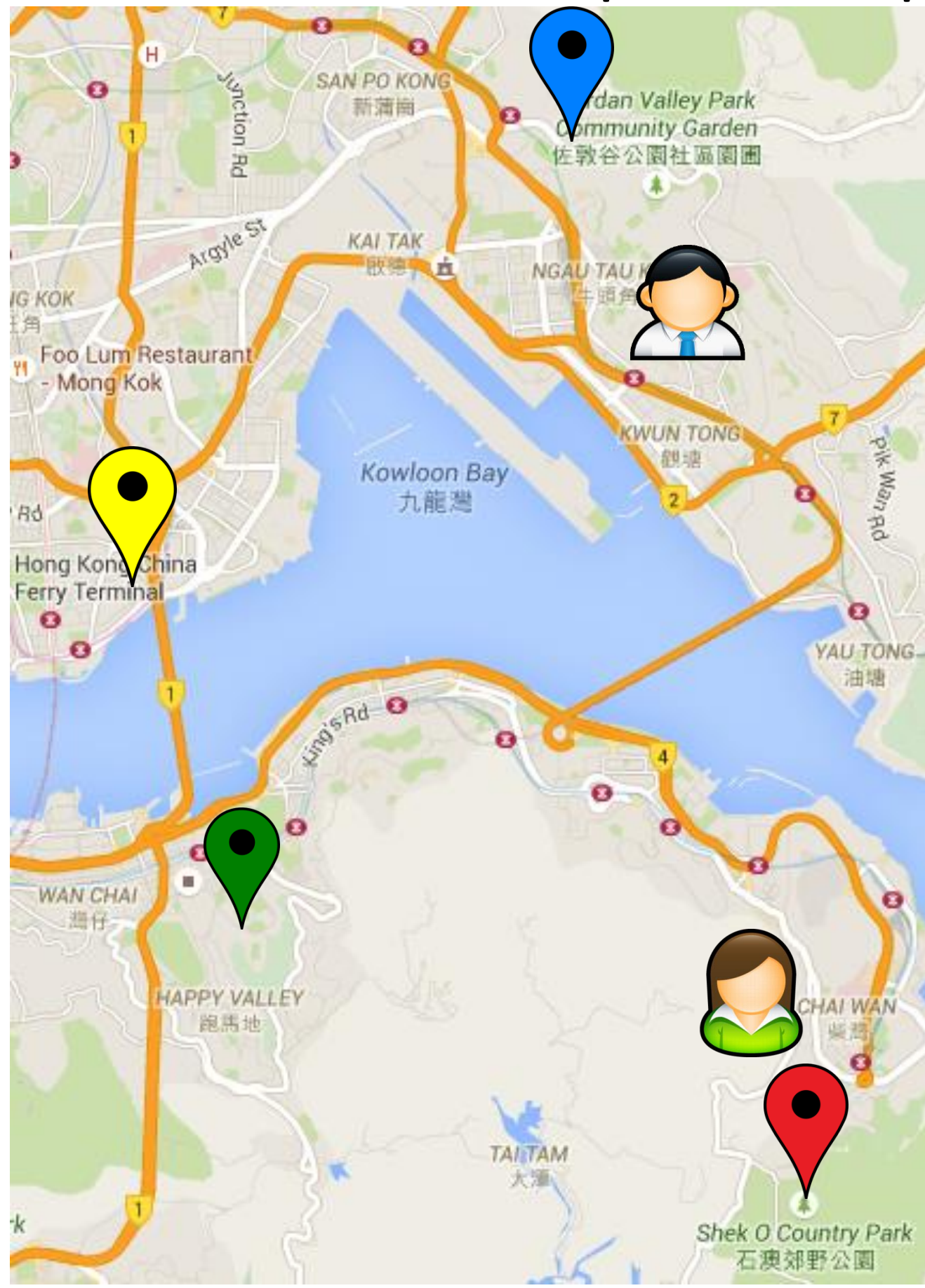
Jieying She[†], Yongxin Tong[‡], Lei Chen[†]

[†]Department of Computer Science and Engineering
The Hong Kong University of Science and Technology
{jshe, leichen}@cse.ust.hk

[‡]State Key Laboratory of Software Development Environment, School of Computer Science and Engineering
Beihang University
yxtong@buaa.edu.cn

Introduction

- Event-Based Social Networks (EBSNs)
 - Online platforms that facilitate offline event organization and participation, e.g. Meetup and Plancast
- Motivation
 - Arrange proper social events to interested users
 - Existing works: either assume user attends one event or ignore location information
 - Spatio-temporal conflicts & travel expenses



- 9:30 AM Hong Kong Hiking Meetup 香港遠足與會
Easy hike (2.2) from Chai Wan MTR to Shek O via Pottinger Peak Country Trail
- 11:00 AM Hong Kong Volunteers 香港義工團
HKKAPS Kindergarten Saturday Morning Playshop From 11 to 12pm
- 1:00 PM Weekends Badminton, Ballroom, Latin Dance & Ad-hoc Dining
23 May, Sat, 1-4 pm, SCAA 12F All Level Games invite 18 players!
- 1:45 PM Hong Kong Dolls, Figures and Toys Collectors Meetup
Lets visit the One Piece 3D Exhibition!

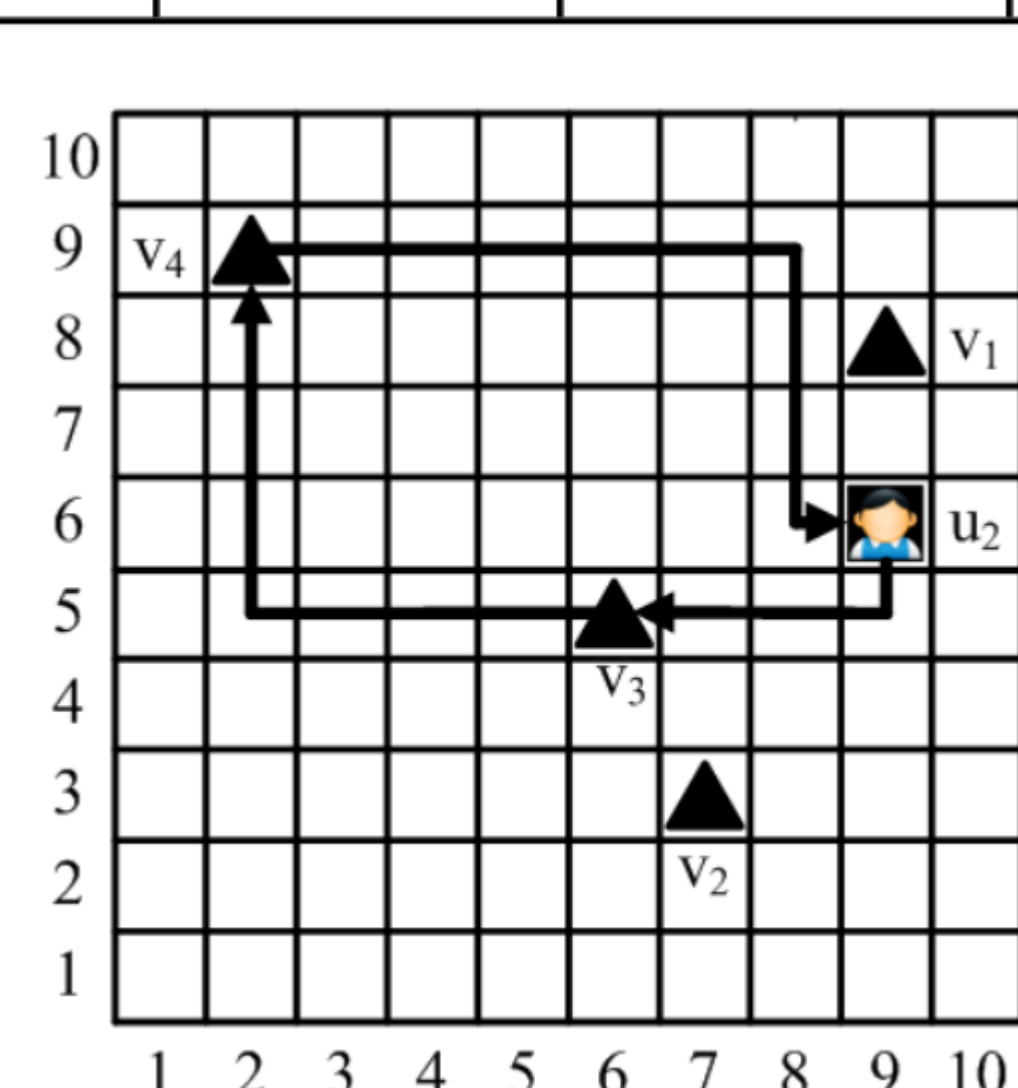
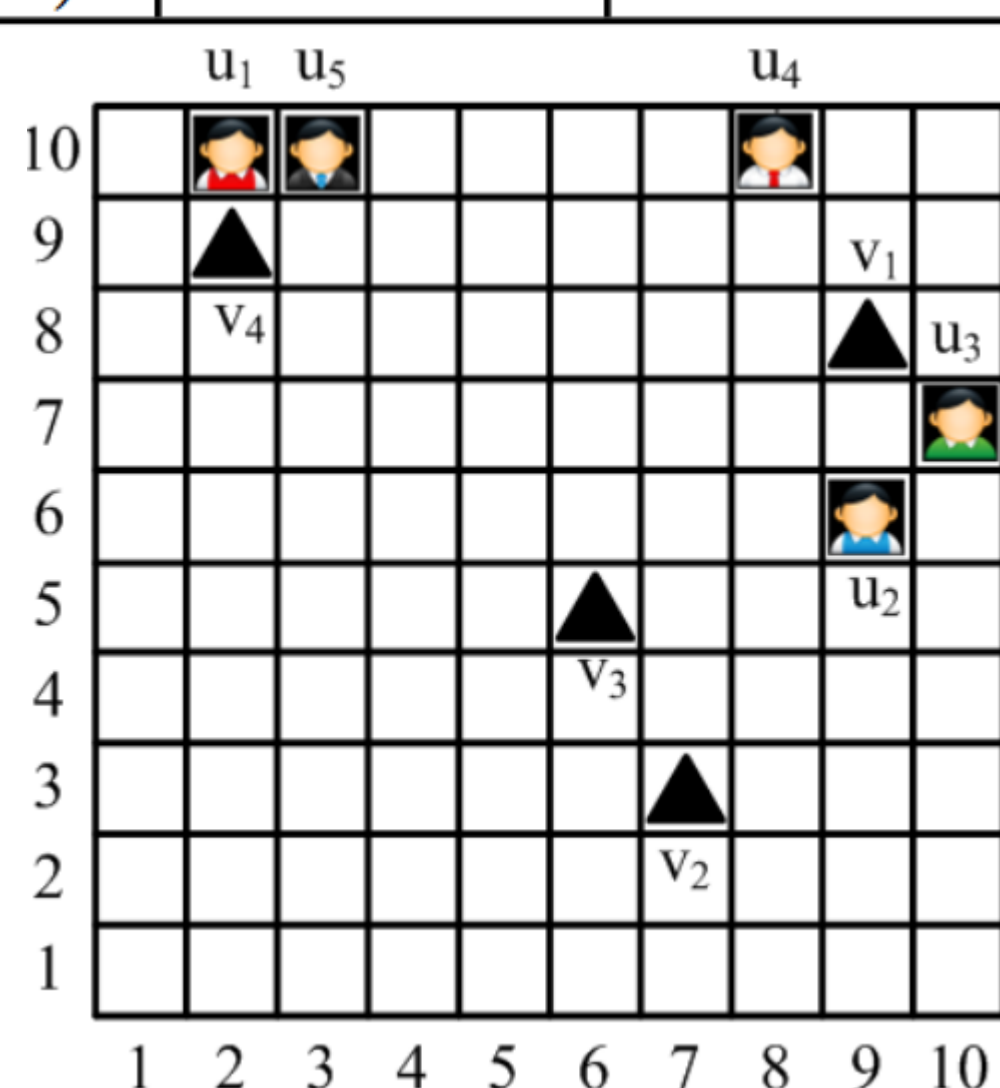
Events on Meetup

The USEP Problem

- Given
 - A set of events V
 - Each $v \in V$: capacity c_v , location l_v , time interval $[t_1^v, t_2^v]$
 - A set of users U
 - Each $u \in U$: location l_u , travel budget b_u
 - Travel cost $\{cost(u, v)\}, \{cost(v_i, v_j)\}$
 - Utility value $\{\mu(v, u)\}$
- Find a planning of schedules $A = \bigcup_u \{S_u\}$
 - Maximize $\Omega(A) = \sum_u \sum_{v \in S_u} \mu(v, u)$
 - Capacities of events are not exceeded
 - No schedule has time conflicts
 - $\mu(v, u) > 0, \forall v \in S_u, \forall u$
 - Travel budgets of users are not exceeded

- The USEP problem is NP-hard

	u_1 (59)	u_2 (29)	u_3 (51)	u_4 (9)	u_5 (33)	Time
v_1 (1)	0.2	0.6	0.7	0.3	0.6	1-4p.m.
v_2 (3)	0.5	0.1	0.3	0.9	0.5	3-6p.m.
v_3 (4)	0.6	0.2	0.9	0.4	0.5	1-2p.m.
v_4 (2)	0.4	0.7	0.2	0.5	0.1	6-7p.m.



Greedy-Based Solution: RatioGreedy

- Maintain a heap, pop a pair with largest *ratio* value each time
 - $ratio(v, u) = \frac{\mu(v, u)}{inc_cost(v, u)}$

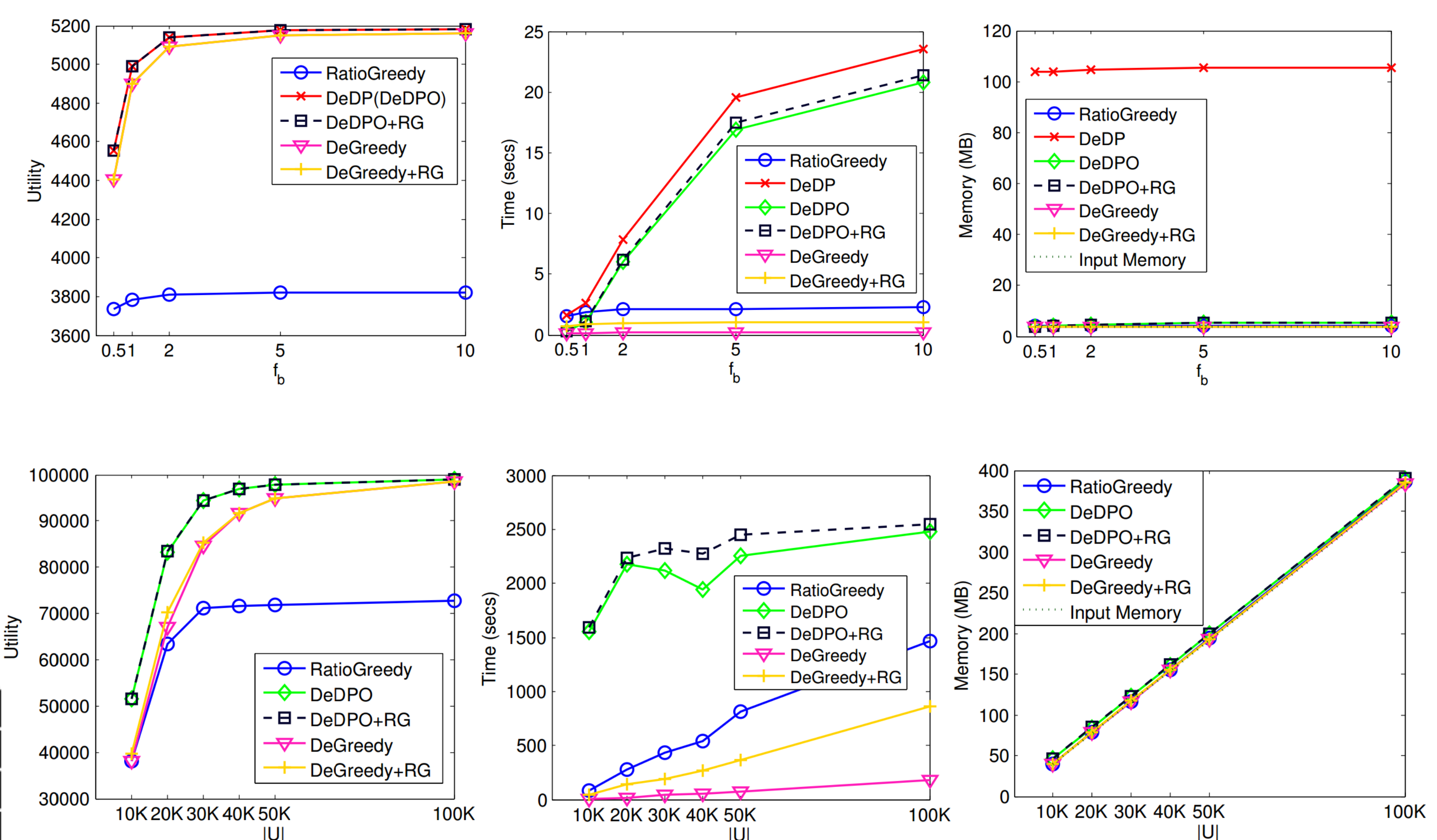
	u_1	u_2	u_3	u_4	u_5
v_1	0.011(18)	0.15(4)	0.175(4)	0.05(6)	0.0375(16)
States of H					
After Initialization	$(v_4, u_1) : 0.2(2), (v_1, u_3) : 0.175(4), (v_1, u_2) : 0.15(4), (v_3, u_3) : 0.075(12), (v_1, u_4) : 0.05(6), (v_1, u_5) : 0.0375(16), (v_2, u_5) : 0.023(22)$				
After 1st Iteration	$(v_1, u_3) : 0.175(4), (v_1, u_2) : 0.15(4), (v_3, u_3) : 0.075(12), (v_1, u_4) : 0.05(6), (v_1, u_5) : 0.0375(16), (v_3, u_1) : 0.0375(16), (v_4, u_2) : 0.035(20), (v_2, u_5) : 0.023(22)$				

Two-Step Approximation Solution: DeDP

- Decomposed into $|U|$ problems
 - Find a schedule for each u with a dynamic programming algorithm
 - Combine the result of each u
- Optimization
 - Optimize space & speed with a proved property
 - Optimize utility with RatioGreedy
- Approximation ratio: $1/2$
- Speed up by replacing DP with a greedy strategy

	$\hat{v}_1$	$\hat{v}_2$	$\hat{v}_3$	$\hat{v}_4$	User	V'_r
u_1	$v_{1,1}(0.2)$	$v_{2,1}(0.5)$	$v_{3,1}(0.6)$	$v_{4,1}(0.4)$	u_1	$\mathbf{v_{3,1}}, v_{1,1}, \mathbf{v_{2,1}}, v_{4,1}$
u_2	$v_{1,1}(0.6)$	$v_{2,2}(0.1)$	$v_{3,2}(0.2)$	$v_{4,1}(0.7)$	u_2	$\mathbf{v_{3,2}}, \mathbf{v_{1,1}}, v_{2,2}, \mathbf{v_{4,1}}$
u_3	$v_{1,1}(0.1)$	$v_{2,2}(0.3)$	$v_{3,2}(0.9)$	$v_{4,2}(0.2)$	u_3	$\mathbf{v_{3,2}}, v_{1,1}, \mathbf{v_{2,2}}, v_{4,2}$
u_4	$v_{1,1}(-0.3)$	$v_{2,3}(0.9)$	$v_{3,3}(0.4)$	$v_{4,2}(0.5)$	u_4	$\emptyset$
u_5	$v_{1,1}(0)$	$v_{2,3}(0.5)$	$v_{3,3}(0.5)$	$v_{4,2}(0.1)$	u_5	$\mathbf{v_{3,3}}, \mathbf{v_{2,3}}, v_{4,2}$

Evaluation



Acknowledgements

This work is supported in part by the Hong Kong RGC Project N_HKUST637/13, National Grand Fundamental Research 973 Program of China under Grant 2014CB340303, NSFC Grant No. 61232018/61300031, Microsoft Research Asia Gift Grant, Google Faculty Award 2013, and Microsoft Research Asia Fellowship 2012.